

# Conjugate Direction Methods

## GIAN Short Course on Optimization: Applications, Algorithms, and Computation

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September 12-24, 2016

# Outline

- 1 Conjugate Direction Methods
- 2 Classical Conjugate Gradient Method
- 3 The Barzilai-Borwein Method



## Exact Line-Search for Quadratics

Analysis uses exact line-search arguments.  
Consider quadratic

$$q(x) = \frac{1}{2}x^T Gx + b^T x$$

and perform an **exact line-search**:  $\hat{x} + \alpha s$ :

$$\underset{\alpha \geq 0}{\text{minimize}} \quad q(\hat{x} + \alpha s) = \frac{1}{2}(\hat{x} + \alpha s)^T G(\hat{x} + \alpha s) + b^T(\hat{x} + \alpha s)$$

Re-arrange quadratic as

$$q(\hat{x} + \alpha s) = \frac{1}{2}\alpha^2 s^T Gs + \alpha (s^T G\hat{x} + b^T s) + \frac{1}{2}\hat{x}^T G\hat{x} + b^T \hat{x}$$

Setting  $\frac{dq}{d\alpha} = 0$  we get:

$$0 = \alpha s^T Gs + s^T (G\hat{x} + b) \quad \Leftrightarrow \quad \alpha = -\frac{s^T (G\hat{x} + b)}{s^T Gs} = \frac{-s^T \nabla q(\hat{x})}{s^T Gs}$$



# Conjugate Direction Methods

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x)$$

Conjugate direction methods relate to a quadratic model of  $f(x)$ .

## Definition (Conjugacy)

$m \leq n$  nonzero vectors,  $s^{(1)}, \dots, s^{(m)} \in \mathbb{R}^n$  are *conjugate wrt positive definite Hessian  $G$* , iff  $s^{(i)T} G s^{(j)} = 0$  for all  $i \neq j$ .

- Conjugacy is orthogonality across positive definite Hessian,  $G$ .
- For  $G = I$ , get orthogonality.

## Definition (Conjugacy)

A *conjugate direction method* generates conjugate directions applied to a positive definite quadratic.



## Conjugate Direction Methods

### Theorem (Linear Independence of Conjugate Directions)

*A set of  $m$  conjugate directions is linearly independent.*

**Proof.**  $s^{(1)}, \dots, s^{(m)} \in \mathbb{R}^n$  conjugate. Consider  $\sum_{i=1}^m a_i s^{(i)} = 0$

... need to show  $a_i = 0$  is only solution of this system

$G$  positive definite  $\Rightarrow G$  nonsingular, hence

$$\sum_{i=1}^m a_i s^{(i)} = 0 \quad \Leftrightarrow \quad G \left( \sum_{i=1}^m a_i s^{(i)} \right) = 0.$$

Pre-multiply by  $s^{(j)}$ :

$$s^{(j)T} G \left( \sum_{i=1}^m a_i s^{(i)} \right) = 0 \quad \Leftrightarrow \quad a_j s^{(j)T} G s^{(j)} = 0 \quad \Leftrightarrow \quad a_j = 0,$$

because  $G$  positive definite. □

# Conjugate Direction Methods

## Theorem (Termination of Conjugate Direction Methods)

- A conjugate direction method terminates for a positive definite quadratic in at most  $n$  exact line-searches.
- Each iterate,  $x^{(k+1)}$  reached by  $k \leq n$  descent steps along conjugate directions  $s^{(1)}, \dots, s^{(k)} \in \mathbb{R}^n$ .

**Proof.** Define the quadratic as

$$q(x) = \frac{1}{2}x^T Gx + b^T x.$$

Conjugate direction,  $s^{(k)}$ , gives  $k + 1$  iterate as

$$x^{(k+1)} = x^{(k)} + \alpha_k s^{(k)} = \dots = x^{(1)} + \sum_{j=1}^k \alpha_j s^{(j)} = x^{(i+1)} + \sum_{j=i+1}^k \alpha_j s^{(j)}.$$



## Conjugate Direction Methods

### Proof cont.

From previous page: Conjugate direction,  $s^{(k)}$ , give iterates

$$x^{(k+1)} = x^{(k)} + \alpha_k s^{(k)} = \dots = x^{(1)} + \sum_{j=1}^k \alpha_j s^{(j)} = x^{(i+1)} + \sum_{j=i+1}^k \alpha_j s^{(j)}.$$

Corresponding gradient of quadratic is

$$\begin{aligned} g^{(k+1)} &= Gx^{(k+1)} + b = G \left( x^{(i+1)} + \sum_{j=i+1}^k \alpha_j s^{(j)} \right) + b \\ &\Rightarrow g^{(k+1)} = g^{(i+1)} + \sum_{j=i+1}^k \alpha_j Gs^{(j)} \end{aligned}$$

Pre-multiply by  $s^{(i)}$  we get

$$s^{(i)T} g^{(k+1)} = s^{(i)T} g^{(i+1)} + \sum_{j=i+1}^k \alpha_j s^{(i)T} Gs^{(j)} = 0, \quad \forall i = 1, \dots, k-1,$$



## Conjugate Direction Methods

### Proof cont.

From previous: pre-multiply by  $s^{(i)}$  we get

$$s^{(i)T} g^{(k+1)} = s^{(i)T} g^{(i+1)} + \sum_{j=i+1}^k \alpha_j s^{(i)T} Gs^{(j)} = 0, \quad \forall i = 1, \dots, k-1,$$

where

- $s^{(i)T} g^{(i+1)} = 0$  due to exact line search.
- $s^{(i)T} Gs^{(j)} = 0$  due to conjugacy.
- $s^{(k)T} g^{(k+1)} = 0$  due to exact line-search.

Hence,

$$s^{(i)T} g^{(k+1)} = 0, \quad \forall i = 1, \dots, k.$$

Now, let  $k = n$ , then it follows that

$$s^{(i)T} g^{(n+1)} = 0, \quad \forall i = 1, \dots, n \quad \Rightarrow \quad g^{(n+1)} = 0$$

because,  $g^{(n+1)}$  orthogonal to  $n$  linearly independent vectors □





# Conjugate Direction Methods

## Remark

*Previous Theorem holds for all conjugate direction methods!*

Methods differ how  $s^{(k)}$  constructed **without knowing Hessian**

## Conjugate Direction Line-Search Method

Given  $x^{(0)}$ , set  $k = 0$ . **repeat**

    Compute the conjugate direction  $s^{(k)}$ .

    Compute the steplength  $\alpha_k := \text{Armijo}(f(x), x^{(k)}, s^{(k)})$

    Set  $x^{(k+1)} := x^{(k)} + \alpha_k s^{(k)}$  and  $k = k + 1$ .

**until**  $x^{(k)}$  is (local) optimum;

... next consider different ways to create conjugate directions.



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- 2 Classical Conjugate Gradient Method
- 3 The Barzilai-Borwein Method



# Classical Conjugate Gradient Method

## Idea Behind Conjugate Gradients

Modify steepest descent so that directions are conjugate.

Start by deriving method for quadratic

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad q(x) = \frac{1}{2}x^T Gx + b^T x$$

then generalize to nonlinear  $f(x)$ .

Start with  $s^{(0)} = -g^{(0)}$ , steepest descent direction

$\Rightarrow$  first step guaranteed to be downhill ... no stalling like Newton!



# Classical Conjugate Gradient Method

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad q(x) = \frac{1}{2}x^T Gx + b^T x$$

Start with  $s^{(0)} = -g^{(0)}$ , steepest descend direction

Choose  $s^{(1)}$  as component of  $-g^{(1)}$  conjugate to  $s^{(0)}$ :

$$s^{(1)} = -g^{(1)} + \beta_0 s^{(0)}$$

Look for formula for  $\beta_0$  such that conjugacy holds, i.e.

$$0 = s^{(0)T} Gs^{(1)} = s^{(0)T} G \left( -g^{(1)} + \beta_0 s^{(0)} \right).$$

Solve for  $\beta_0$ , and get

$$\beta_0 = \frac{s^{(0)T} Gg^{(1)}}{s^{(0)T} Gs^{(0)}},$$

where  $s^{(0)T} Gs^{(0)} \neq 0$ , because  $G$  positive definite, and  $s^{(0)} \neq 0$ .



# Classical Conjugate Gradient Method

Simplify formula for  $\beta_0$ :

$$\beta_0 = \frac{s^{(0)T} Gg^{(1)}}{s^{(0)T} Gs^{(0)}},$$

Recall, that

$$x^{(1)} = x^{(0)} + \alpha_1 s^{(0)} \Leftrightarrow s^{(0)} = (x^{(1)} - x^{(0)}) / \alpha_1,$$

where  $\alpha_1 \neq 0$ , because of steepest descend.

Now use  $G\delta = \gamma$  to write  $\beta_0$  as

$$\beta_0 = \frac{(x^{(1)} - x^{(0)})^T Gg^{(1)}}{(x^{(1)} - x^{(0)})^T Gs^{(0)}} = \frac{(g^{(1)} - g^{(0)})^T g^{(1)}}{(g^{(1)} - g^{(0)})^T s^{(0)}}$$

Exact line-search implies  $0 = g^{(1)T} s^{(0)} = -g^{(1)T} g^{(0)}$ , and thus

$$\beta_0 = \frac{g^{(1)T} g^{(1)}}{g^{(0)T} g^{(0)}}.$$



# Classical Conjugate Gradient Method

Consider general step,  $k$ :

$s^{(k)}$  = the component of  $-g^{(k)}$  conjugate to  $s^{(0)}, \dots, s^{(k-1)}$ .

Desired conjugacy:

$$s^{(k)T} G s^{(j)} = 0, \forall j < k \quad \Leftrightarrow \quad s^{(k)T} \gamma^{(j)} = 0, \forall j < k,$$

Use Gram-Schmidt orthogonalization procedure to get

$$s^{(k)} = -g^{(k)} + \sum_{j=0}^{k-1} \beta_j s^{(j)} \quad \text{Can } \beta_j = 0 \text{ for } j < k???$$

For quadratic, can show that  $\beta_j = 0, \forall j < k$ . Hence get:

$$s^{(k)} = -g^{(k)} + \beta_{k-1} s^{(k-1)} \quad \text{where} \quad \beta_{k-1} = \begin{cases} 0 & \text{if } k = 0 \\ \frac{g^{(k)T} g^{(k)}}{g^{(k-1)T} g^{(k-1)}} & \text{otherwise} \end{cases}$$

Fletcher-Reeves conjugate gradient method

# Classical Conjugate Gradient Method

Min. quadratic  $q(x) = \frac{1}{2}x^T Gx + b^T x$  with Fletcher-Reeves (FR)

$$s^{(k)} = g^{(k)} + \beta_{k-1} s^{(k-1)} \quad \text{where} \quad \beta_k = \begin{cases} 0 & \text{if } k = -1 \\ \frac{g^{(k)T} g^{(k)}}{g^{(k-1)T} g^{(k-1)}} & \text{otherwise} \end{cases}$$

## Theorem (Convergence of FR for Convex Quadratics)

FR with exact line-search terminates at stationary point,  $x^{(m)}$  after  $m \leq n$  iterations for a pos. definite quadratic. Moreover, for  $0 \leq i \leq m - 1$ , we have that:

- 1 Conjugate search directions:  $s^{(i)T} G s^{(j)} = 0 \quad \forall i \neq j, j < i$ .
- 2 Orthogonal gradients:  $g^{(i)T} g^{(j)} = 0 \quad \forall i \neq j, j = 1, \dots, i - 1$ .
- 3 Descend property:  $s^{(i)T} g^{(j)} = -g^{(i)T} g^{(j)} < 0 \quad \forall i \neq j$ .



# Proof of Fletcher-Reeves Convergence

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- 3 Descend property:  $s^{(i)T} g^{(i)} = -g^{(i)T} g^{(i)} < 0 \forall i \neq j$ .

**Proof.** By induction over  $m$  ...

For  $m = 0$ , there is nothing to show.

For  $m \geq 1$ , show 1. to 3. of Theorem by induction over  $i$ .

For  $i = 0$ , observe

$$s^{(0)} = -g^{(0)} \Rightarrow s^{(0)T} g^{(0)} = -g^{(0)T} g^{(0)}.$$

$\Rightarrow$  3. holds for  $i = 0$ , nothing to show for 1. and 2. (no  $j < 0$ !)





# Proof of Fletcher-Reeves Convergence

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- 1 Conjugate search directions:  $s^{(i)T} Gs^{(j)} = 0 \forall i \neq j, j < i$ .
- 2 Orthogonal gradients:  $g^{(i)T} g^{(j)} = 0 \forall i \neq j, j = 1, \dots, i - 1$ .
- 3 Descend property:  $s^{(i)T} g^{(i)} = -g^{(i)T} g^{(i)} < 0 \forall i \neq j$ .

**Proof cont.** Induction hypothesis: Assume that 1.-3. hold for  $i$ . Show 1.-3. also hold for  $i + 1$ : Quadratic objective implies:

$$g^{(i+1)} = Gx^{(i+1)} + b = G \left( x^{(i)} + \alpha_i s^{(i)} \right) + b = g^{(i)} + \alpha_i Gs^{(i)}$$

Exact line search  $\alpha_i$  implies:

$$\alpha_i = \frac{-g^{(i)T} s^{(i)}}{s^{(i)T} Gs^{(i)}} = \frac{g^{(i)T} g^{(i)}}{s^{(i)T} Gs^{(i)}}, \quad \text{from 3. by induction}$$



# Proof of Fletcher-Reeves Convergence

Now, we consider Part 2 for  $g^{(i)T} g^{(j)} = 0$ :

$$\begin{aligned}g^{(i+1)T} g^{(j)} &= g^{(i)T} g^{(j)} + \alpha_i s^{(i)T} G g^{(j)} \\ &= g^{(i)T} g^{(j)} + \alpha_i s^{(i)T} G \left( -s^{(j)} + \beta_{j-1} s^{(j-1)} \right)\end{aligned}$$

from definition of  $s^{(j)} = -g^{(j)} + \beta_{j-1} s^{(j-1)}$ , to get  $g^{(j)}$ . Thus,

$$g^{(i+1)T} g^{(j)} = g^{(i)T} g^{(j)} - \alpha_i s^{(i)T} G s^{(j)} + \alpha_i \beta_{j-1} s^{(i)T} G s^{(j-1)}$$

For  $i = j$  observe:

- Exact line-search  $\Rightarrow \alpha = \frac{-s^T g}{s^T G s} \Rightarrow$  sum of first terms is zero
- Induction Part 1.  $\Rightarrow$  last expression zero.



## Proof of Fletcher-Reeves Convergence

Now, we consider Part 2 for  $g^{(i+1)T} g^{(j)} = 0$ :

$$g^{(i+1)T} g^{(j)} = g^{(i)T} g^{(j)} - \alpha_i s^{(i)T} G s^{(j)} + \alpha_i \beta_{j-1} s^{(i)T} G s^{(j-1)}$$

For  $i < j$  observe:

- Induction Part 2.  $\Rightarrow$  first expression zero
- Induction Part 1.  $\Rightarrow$  last two expressions zero.

Thus,  $g^{(i+1)T} g^{(j)} = 0$  for  $j = 1, \dots, i$  which proves Part 2.



## Proof of Fletcher-Reeves Convergence

Consider Part 1. Use  $s^{(i+1)} = -g^{(i+1)} + \beta_i s^{(i)}$ :

$$\begin{aligned} s^{(i+1)T} Gs^{(j)} &= -g^{(i+1)T} Gs^{(j)} + \beta_i s^{(i)T} Gs^{(j)} \\ &= \alpha_j^{-1} g^{(i+1)T} (g^{(j)} - g^{(j+1)}) + \beta_i s^{(i)T} Gs^{(j)}, \end{aligned}$$

because  $Gs^{(j)} = \alpha_j^{-1} G(x^{(j)} - x^{(j+1)}) = \alpha_j^{-1} G(g^{(j)} - g^{(j+1)})$ .

For  $j < i$  get:

- Part 2.  $\Rightarrow$  first component is zero
- Part 1. and induction  $\Rightarrow$  second component is zero



## Proof of Fletcher-Reeves Convergence

Consider again

$$\begin{aligned} s^{(i+1)T} G_S(j) &= -g^{(i+1)T} G_S(j) + \beta_i s^{(i)T} G_S(j) \\ &= \alpha_j^{-1} g^{(i+1)T} \left( g^{(j)} - g^{(j+1)} \right) + \beta_i s^{(i)T} G_S(j), \end{aligned}$$

For  $j = i$  re-write this expression as

$$s^{(j+1)T} G_S(j) = \alpha_j^{-1} g^{(j+1)T} g^{(j)} - \alpha_j^{-1} g^{(j+1)T} g^{(j+1)} + \beta_j s^{(j+1)T} G_S(j).$$

Part 2.  $\Rightarrow$  first component is zero

Use exact line-search  $\alpha_j$  second component becomes

$$\begin{aligned} & -\alpha_j^{-1} g^{(j+1)T} g^{(j+1)} + \beta_j s^{(j+1)T} G_S(j) \\ &= -s^{(j+1)T} G_S(j) \frac{g^{(j+1)T} g^{(j+1)}}{g^{(j)T} g^{(j)}} + \beta_j s^{(j+1)T} G_S(j) = 0, \end{aligned}$$

from  $\beta_j$  formula.

$\Rightarrow s^{(i+1)T} G_S(j) = 0$  for all  $j = 1, \dots, i$ , which proves Part 1.

Quadratic termination follows from Part 1., and conjugate directions,  $s^{(1)}, \dots, s^{(m)}$ . □



# Conjugate Gradient Method for Nonlinear Functions

Consider minimize  $f(x)$ , then  
 $x \in \mathbb{R}^n$

- Cannot perform exact line-search ... approx, e.g. Armijo
- Cannot expect termination after  $n$  steps  
 $\Rightarrow$  re-start  $s^{(n+1)} = -g^{(n+1)}$  or re-orthogonalize

## Other Conjugate Gradient Schemes

$$\beta_k^{PR} = \frac{(g^{(k+1)} - g^{(k)})^T g^{(k)}}{g^{(k-1)T} g^{(k-1)}}$$

and

$$\beta_k^{DY} = \frac{g^{(k)T} g^{(k)}}{s^{(k-1)T} g^{(k-1)}}$$

Dai-Yuan better than Polak-Ribiere better than Fletcher-Reeves



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# The Barzilai-Borwein Method

Recent renewed interest in a simpler two-step gradient method

- Satisfy quasi-Newton in least-squares sense.

## Barzilai-Borwein Method

Given  $x^{(0)}$ , set  $k = 0$ .

**repeat**

    Set the step-size  $\alpha_k$  using one of BB schemes below.

    Set  $x^{(k+1)} := x^{(k)} - \alpha_k g^{(k)}$  and  $k = k + 1$ . [Steepest Descent]

**until**  $x^{(k)}$  is (local) optimum;

## Surprise: No Line Search

- Barzilai-Borwein Algorithm has **no** line-search
- Success relies on non-monotone behavior (may increase  $f(x)$ )





# The Barzilai-Borwein Method

Popular formulas for BB step size

$$\alpha_k^{BB} = \frac{\delta^{(k-1)}\delta^{(k-1)}}{\delta^{(k-1)}\gamma^{(k-1)}} \quad (1)$$

$$\alpha_k^{BBs} = \frac{\delta^{(k-1)}\gamma^{(k-1)}}{\gamma^{(k-1)}\gamma^{(k-1)}} \quad (2)$$

$$\alpha_k^{aBB} = \begin{cases} \alpha_k^{BB} & \text{for odd } k \\ \alpha_k^{BBs} & \text{for even } k \end{cases} \quad (3)$$

- Can reset the step length to steepest-descend
- Generalized to bound-constrained optimization using projection



# Summary of Conjugate Direction Methods

Methods for unconstrained optimization:

$$\underset{x}{\text{minimize}} f(x)$$

- Conjugacy is orthogonality across Hessian  $G$ , i.e.

$$s^{(i)T} G s^{(j)} = 0 \quad \forall i \neq j$$

- Conjugate direction methods terminate finitely for quadratic
- Good alternative to quasi-Newton
- Recently, interest in Barzilai-Borwein schemes

